

Optimizing MPI Performance on Theta Using MPI on ThetaGPU

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Sudheer Chunduri sudheer@anl.gov



Outline

- Theta
 - Cray XC network software stack and MPI software stack
 - Non-blocking collectives
 - Topology mapping optimizations
 - Few key performance tuning knobs
 - Cray XC routing optimizations
 - Darshan-Autoperf
- ThetaGPU
 - Introduction to MPI usage on ThetaGPU

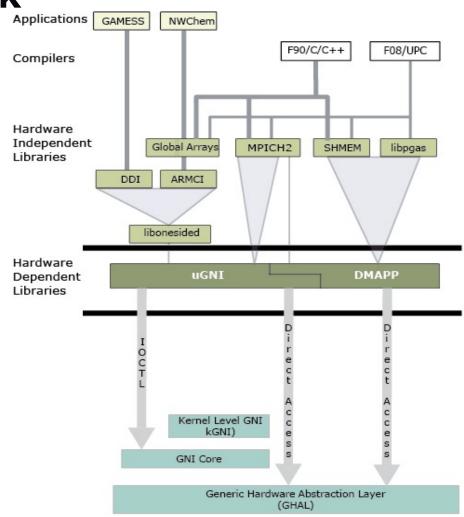


Cray XC Network Software Stack

uGNI - Generic Network Interface (message passing based)

DMAPP - Distributed Shared Memory Application APIs (shared memory)

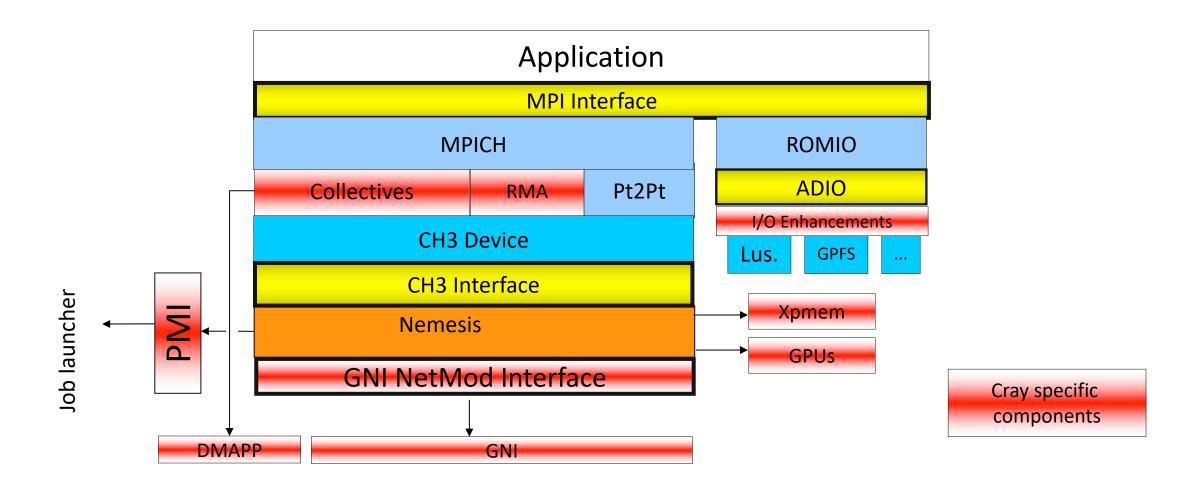
uGNI and **DMAPP** provide low-level communication services to user-space software



Hardware



Cray MPI Software Stack (CH3 device)





MPI-3 Nonblocking Collectives

- Enables overlap of communication/computation similar to nonblocking (send/recv) communication
- Non-blocking variants of all collectives: MPI_lbcast (<bcast args>, MPI_Request *req);
- Semantics
 - Function returns no matter what
 - Usual completion calls (wait, test)
 - Out-of-order completion
- Semantic advantages
 - Enables asynchronous progression (software pipelining)
 - Decouple data transfer and synchronization (Noise Resiliency)
 - Allow overlapping communicators
 - Multiple outstanding operations at any time

```
MPI_Comm comm;
int array1[100], array2[100];
int root=0;
MPI_Request req;
...
MPI_Ibcast(array1, 100, MPI_INT,
root, comm, &req);
compute(array2, 100);
MPI_Wait(&req, MPI_STATUS_IGNORE);
```

MPI-3 Nonblocking Collectives Support

- Includes many optimizations for MPI-3 nonblocking Collectives
- Not ON by default. User must set the following env. Variables:

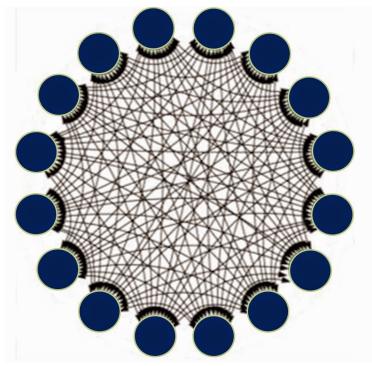
 export MPICH_NEMESIS_ASYNC_PROGRESS=[SC|MC|ML] (network interface DMA engine enables asynchronous progress)
 export MPICH_MAX_THREAD_SAFETY=multiple
- Special optimizations for Small message MPI_Iallreduce, based on Aries HW Collective Engine:

```
Users must link against DMAPP
-Wl,--whole-archive,-ldmapp,--no-whole-archive (static linking)
-ldmapp (dynamic linking)
export MPICH_NEMESIS_ASYNC_PROGRESS = [SC|MC|ML]
export MPICH_MAX_THREAD_SAFETY=multiple
export MPICH_USE_DMAPP_COLL=1
```



Topology Mapping and Rank Reordering

- Topology mapping
 - Minimize communication costs through interconnect topology aware task mapping
 - Could *potentially* help reduce congestion
 - Node placement for the job could be a factor (no explicit control available to request a specific placement)
- Application communication pattern
 - MPI process topologies expose this in a portable way
 - Network topology agnostic
- Rank reordering
 - Can override the default mapping scheme
 - The default policy for **aprun** launcher is SMP-style placement
 - To display the MPI rank placement information,
 - set MPICH_RANK_REORDER_DISPLAY.





MPI Rank Reordering

- MPICH_RANK_REORDER_METHOD
 - Vary rank placement to optimize communication (ex: maximize on-node communication between MPI ranks)
 - Use CrayPat with "-g mpi" to produce a specific MPICH_RANK_ORDER file to maximize intra-node communication
 - Or, use perf_tools grid_order command with your application's grid dimensions to layout
 MPI ranks in alignment with data grid
 - To use:
 - name your custom rank order file: MPICH_RANK_ORDER
 - This approach is physical system topology agnostic
 export MPICH_RANK_REORDER_METHOD=3



MPI Rank Reordering

- MPICH_RANK_REORDER_METHOD (cont.)
 - A topology and placement aware reordering method is also available
 - Optimizes rank ordering for Cartesian decompositions using the layout of nodes in the job
 - To use:
 - export MPICH_RANK_REORDER_METHOD=4
 - export MPICH_RANK_REORDER_OPTS="-ndims=3 -dims=16,16,8"

MPI Grid Detection:

There appears to be point-to-point MPI communication in a 96 X 8 grid pattern. The 52% of the total execution time spent in MPI functions might be reduced with a rank order that maximizes communication between ranks on the same node. The effect of several rank orders is estimated below.

A file named MPICH_RANK_ORDER.Grid was generated along with this report and contains usage instructions and the Custom rank order from the following table.

		MPICH_RANK_REORDER_METHOD
2.385e+09	95.55%	3
1.880e+09	75.30%	1
1.373e+06	0.06%	2
n 0.000e+00	0.00%	0
	Bytes/PE of Total Bytes/PE 2.385e+09 1.880e+09 1.373e+06	Bytes/PE Bytes/PE% of Total



Profiling with CrayPat

- Application built with
 "pat_build -g mpi"
- pat_report generates the CrayPat report
- Note the MPI call times and number of calls
- Load imbalance across the ranks

Table 1: Profile by Function Group and Function
Time% Time Imb. Time Imb. Calls Group
Time% Function
100.0% 667.935156 49,955,946.2 Total
40.0% 267.180169 49,798,359.2 MPI
24.0 % 160.400193 28.907525 15.3% 2,606,756.0 MPI_Wait
6.4% 42.897564 0.526996 1.2% 157,477.0 MPI_Allreduce
4.8% 31.749303 3.923541 11.0% 42,853,974.0 MPI_Comm_rank
3.5% 23.303805 1.774076 7.1% 1,303,378.0 MPI_Isend
1.1% 7.658009 0.637044 7.7% 1,303,378.0 MPI_Irecv
39.1% 260.882504 2.0 USER
39.1% 260.882424 17.270557 6.2% 1.0 main
20.9% 139.872482 157,585.0 MPI_SYNC
20.4 % 136.485384 36.223589 26.5% 157,477.0 MPI_Allreduce(sync)



Profiling with CrayPat

MPI message sizes are reported

The message size distributions can help characterize an application as

- Latency sensitive
- Bandwidth sensitive

Total		========
MPI Msg Bytes%	100.0%	
MPI Msg Bytes	18,052,938,280.0	
MPI Msg Count	1,460,959.0	msgs
MsgSz <16 Count	157,529.0	msgs
16<= MsgSz <256 Count	65.0	msgs
256<= MsgSz <4KiB Count	2,815.0	msgs
4KiB<= MsgSz <64KiB Count	1,300,511.0	msgs
64KiB<= MsgSz <1MiB Count	39.0	msgs
MPI_Isend		
MPI_Isend		
MPI Msg Bytes%		msgs
MPI Msg Bytes% MPI Msg Bytes	18,051,670,432.0 1,303,378.0	msgs
MPI Msg Bytes% MPI Msg Bytes MPI Msg Count	18,051,670,432.0 1,303,378.0 16.0	=
MPI Msg Bytes% MPI Msg Bytes MPI Msg Bytes MPI Msg Count MsgSz <16 Count	18,051,670,432.0 1,303,378.0 16.0	msgs msgs
MPI Msg Bytes% MPI Msg Bytes MPI Msg Count MsgSz <16 Count 16<= MsgSz <256 Count	18,051,670,432.0 1,303,378.0 16.0 0.0	msgs msgs
MPI Msg Bytes% MPI Msg Bytes MPI Msg Count MsgSz <16 Count 16<= MsgSz <256 Count 256<= MsgSz <4KiB Count	18,051,670,432.0 1,303,378.0 16.0 0.0 2,812.0 1,300,511.0	msgs msgs



Key Environment Variables for XC

- Use MPICH USE DMAPP COLL for hardware supported collectives
 - Most of MPI's optimizations are enabled by default, but not the DMAPP-optimized features, because...
 - Using DMAPP may have some disadvantages
 - May reduce resources MPICH has available (share with DMAPP)
 - Requires more memory (DMAPP internals)
 - DMAPP does not handle transient network errors
 - These are highly-optimized algorithms which may result in significant performance gains, but user has to request them
 - Supported DMAPP-optimized functions:
 - MPI_Allreduce (4-8 bytes)
 - MPI_Bcast (4 or 8 bytes)
 - MPI_Barrier
 - To use (link with libdmapp):
 - Collective use: export MPICH_USE_DMAPP_COLL=1

Key Environment Variables for XC

- MPICH GNI environment variables
 - To optimize inter-node traffic using the Aries interconnect, the following are the most significant env variables to play with (avoid significant deviations from the default if possible):
 - MPICH GNI MAX VSHORT MSG SIZE
 - Controls max message size for E0 mailbox path (Default: varies)
 - MPICH GNI MAX EAGER MSG SIZE
 - Controls max message size for E1 Eager Path (Default: 8K bytes)
 - MPICH GNI NUM BUFS
 - Controls number of 32KB internal buffers for E1 path (Default: 64)
 - MPICH GNI NDREG MAXSIZE
 - Controls max message size for R0 Rendezvous Path (Default: 4MB)
 - MPICH GNI RDMA THRESHOLD
 - Controls threshold for switching to BTE from FMA (Default: 1K bytes)
- Refer the MPI man page for further details

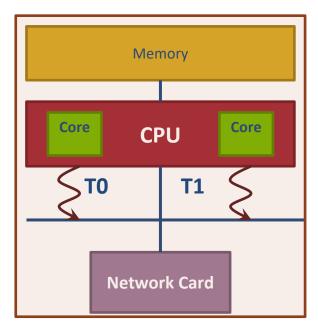
Key Environment Variables for XC

- Specific Collective Algorithm Tuning
 - Different algorithms may be used for different message sizes in collectives (e.g.)
 - Algorithm A might be used for Alltoall for messages < 1K.
 - Algorithm B might be used for messages >= 1K.
 - To optimize a collective, you can modify the cutoff points when different algorithms are used. This may improve performance. A few important ones are:
 - MPICH_ALLGATHER_VSHORT_MSG
 - MPICH_ALLGATHERV_VSHORT_MSG
 - MPICH_GATHERV_SHORT_MSG
 - MPICH_SCATTERV_SHORT_MSG
 - MPICH_GNI_A2A_BLK_SIZE
 - MPICH_GNI_A2A_BTE_THRESHOLD
- Refer the MPI man page for further details



MPI+X Hybrid Programming Optimizations

- MPI Thread Multiple Support for
 - Point to point operations & Collectives (optimized global lock)
 - MPI-RMA (thread hot)
- All supported in default library
 (Non-default Fine-Grained Multi-Threading library is no longer needed)
- Users must set the following env. variable:
 - export MPICH_MAX_THREAD_SAFETY=multiple
- Global lock optimization ON by default (N/A for MPI-RMA)
 - export MPICH_OPT_THREAD_SYNC=0 falls back to pthread_mutex()
- "Thread hot" optimizations for MPI-3 RMA:
 - Contention free progress and completion
 - High bandwidth and high message rate
 - Independent progress thread(s) flush outstanding traffic, other threads make uninterrupted progress

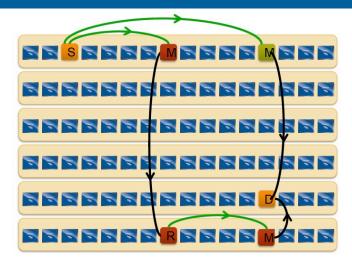


MPI + Threads



Cray XC Routing

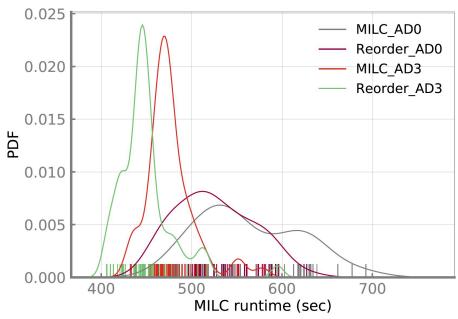
- Aries provides three basic routing modes
 - Deterministic (minimal)
 - Hashed deterministic (minimal, non-minimal), hash on "address"
 - Adaptive
 - 0 No bias (default)
 - 1 Increasing bias towards minimal (as packet travels)
 - Used for MPI all-to-all
 - 2 Straight minimal bias (non-increasing)
 - 3 Strong minimal bias (non-increasing)
- Non-adaptive modes are more susceptible to congestion unless the traffic is very uniform and well-behaved
- MPICH_GNI_ROUTING_MODE environment variable
 - Set to one of ADAPTIVE_[0123], MIN_HASH, NMIN_HASH, IN_ORDER
 - MPICH GNI A2A ROUTING MODE also available



Cray XC group:

Minimal path: 2 hops

Non-minimal path: 4 hops



module load adaptive-routing-a3

module unload adaptive-routing-a3 module help adaptive-routing-a3

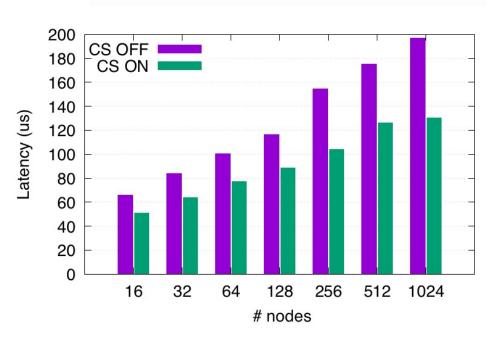


Core Specialization

- Offloads some kernel and MPI work to unused Hyper-Thread(s)
- Good for large jobs and latency sensitive MPI collectives
- Highest numbered unused thread on node is chosen
 - Usually the highest numbered HT on the highest numbered physical core
- Examples
 - aprun -r 1 ...
 - aprun -r N ... # use extra threads
- Cannot oversubscribe, OS will catch
 - Illegal: aprun -r1 -n 256 -N 256 -j 4 a.out
 - Legal: aprun -r1 -n 255 -N 255 -j 4 a.out
 - Legal: aprun -r8 -n 248 -N 248 -j 4 a.out

8-Byte Allreduce on Theta

64 processes per node (run in production — other jobs are running)





Autoperf – Low-overhead MPI profiler

- Autoperf 2.0
 - Collection of modules for profiling MPI (APMPI), network counters (APXC) etc.
 - Based on PMPI (MPI profiling interface)
 - Intercepts 74 MPI operations (collectives, P2P (blocking and non-blocking), RMA)
 - For each operation,
 - Call count
 - Total time per call
 - Message size per op (have 6 bins 1B-256B, 256—1KB, 1KB-8KB, 8KB-256KB, 256-1MB, 1MB+)
 - Sync enable option for collectives adds an extra Barrier
 - By adding this overhead, we can have an idea of load imbalance across the ranks
 - Autoperf is now a submodule within Darshan
 - module load darshan/3.3.0 (on Theta and on ThetaGPU, a new Darshan release with Autoperf included)
 - APXC Another module to capture network counters
 - APNVGPU work in progress to capture GPU performance counters



Autoperf 2.0 output sample

```
\# *******************
# APMPI module data
# description of APMPI counters:
# global summary stats showing the variance across all MPI processes:
    MPI TOTAL COMM TIME VARIANCE: variance in total communication time across all the
processes.
    MPI TOTAL COMM SYNC TIME VARIANCE: variance in total sync time across all the
processes, if enabled.
  per-process detailed stats on the usage of various MPI routines:
    MPI PROCESSOR NAME: name of the processor used by the MPI process.
    MPI * CALL COUNT: total call count for an MPI op.
    MPI * TOTAL BYTES: total bytes (i.e., cumulative across all calls) moved with an MPI op.
    MPI * MSG SIZE AGG *: histogram of total bytes moved for all the calls of an MPI op.
    MPI_*_TOTAL_TIME: total time (i.e, cumulative across all calls) of an MPI op.
    MPI * MIN TIME: minimum time across all calls of an MPI op.
    MPI * MAX TIME: maximum time across all calls of an MPI op.
   MPI * TOTAL SYNC TIME: total sync time (cumulative across all calls of an op) of an MPI op, if
enabled.
    MPI TOTAL COMM TIME: total communication (MPI) time of a process across all the MPI ops.
```

MPI TOTAL COMM SYNC TIME: total sync time of a process across all the MPI ops, if enabled.

```
APMPI 0 7713204788897956166 MPI TOTAL COMM TIME VARIANCE 0.000106
APMPI 0 2250748219237254214 MPI PROCESSOR NAME nid00000
APMPI 0 2250748219237254214 MPI SEND CALL COUNT 2100
APMPI 0 2250748219237254214 MPI SEND TOTAL BYTES 734006000
APMPI 0 2250748219237254214 MPI_SEND_MSG_SIZE_AGG_0_256 700
APMPI 0 2250748219237254214 MPI SEND MSG SIZE AGG 256 1K 0
APMPI 0 2250748219237254214 MPI SEND MSG SIZE AGG IK 8K 0
APMPI 0 2250748219237254214 MPI_SEND_MSG_SIZE_AGG_8K_256K 0
APMPI 0 2250748219237254214 MPI SEND MSG SIZE AGG 256K IM 1400
APMPI 0 2250748219237254214 MPI SEND MSG SIZE AGG IM PLUS 0
APMPI 0 2250748219237254214 MPI SEND TOTAL TIME 0.207326
APMPI 0 2250748219237254214 MPI SEND MIN TIME 0.000003
APMPI 0 2250748219237254214 MPI SEND MAX TIME 0.002291
```



MPI on Theta GPU



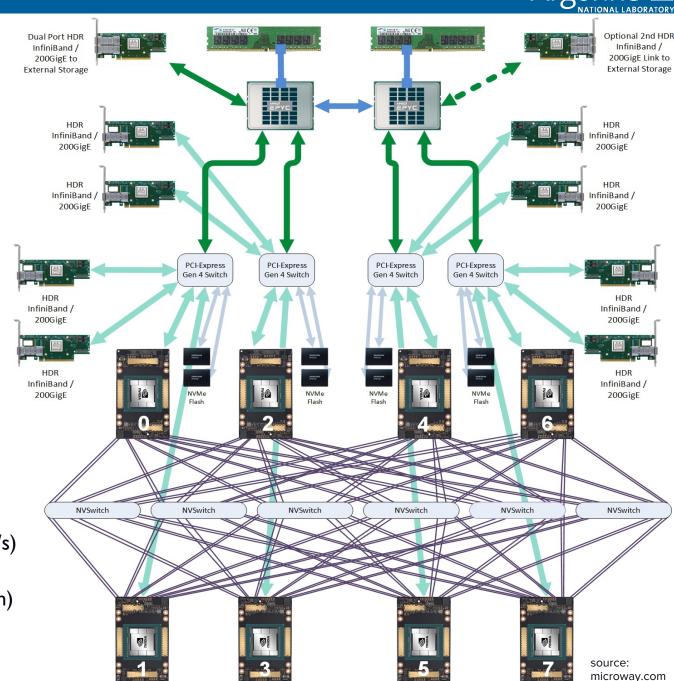
ThetaGPU node

- NVIDIA DGX A100
- 2 AMD EPYC Rome CPUs
- 8 NVIDIA A100 GPUs

AMD Infinity fabric

- GPUs interconnected over NVLink
- 8 HDR IB NICs
- PCIe Gen 4 switch for connecting CPUs, GPUs and NICs

Eight-channel DDR4 (171GB/s per CPU)
PCI-Express x16 (gen 4.0) bus to CPU
PCI-Express x16 (gen 4.0) GPUDirectRDMA (~24.7 GB/s)
PCI-Express x8 (gen 4.0) GPUDirectStorage
25GB/s NVLink 3.0(50GB/s bi-directional per connection)
(2 connections from GPU to each of 6 NVSwitches
600 GB/s of NVLink BW for each GPU)



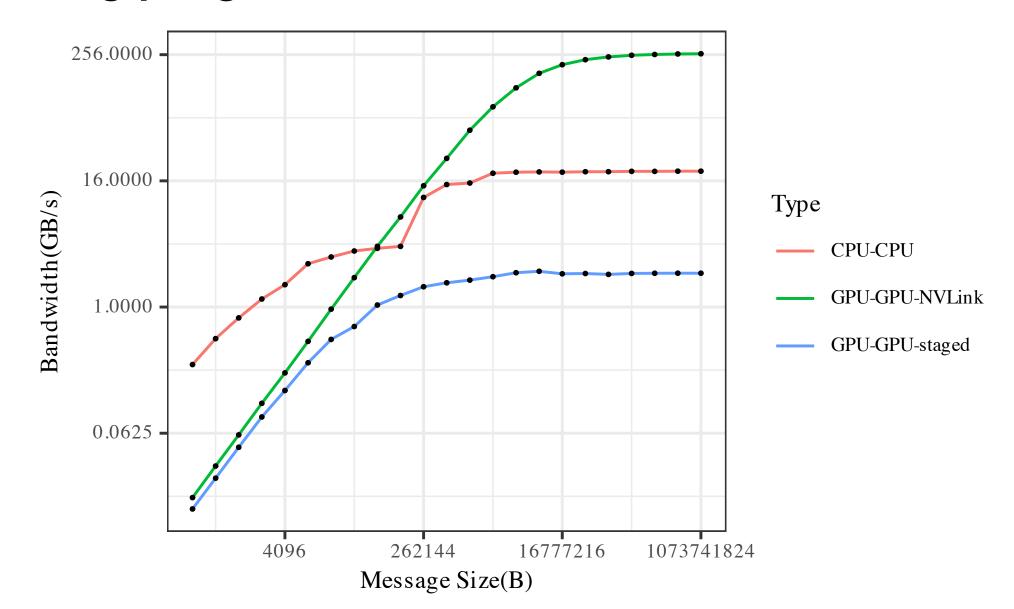


MPI Usage on ThetaGPU

- OpenMPI over UCX is the default MPI currently
- MPICH over OFI will be installed as well soon.
- Current working recipe:
 - Compilation:
 - /soft/thetagpu/hpc-sdk/Linux_x86_64/21.3/cuda/bin/nvcc -l/soft/thetagpu/hpc-sdk/Linux_x86_64/21.3/comm_libs/mpi/include -c test.cu
 - /soft/thetagpu/hpc-sdk/Linux_x86_64/21.3/cuda/bin/nvcc -L/soft/thetagpu/hpc-sdk/Linux_x86_64/21.3/comm_libs/mpi/lib -lmpi -o test test.o
 - Run:
 - mpirun -n 2 -bind-to socket -x UCX_TLS=cuda_ipc ./test
 - Example:
 - https://github.com/argonne-lcf/CompPerfWorkshop-2021/tree/main/07 mpi/MPI ping-pong
 - Compares the ping pong latency and bandwidth for message sizes (8B IMB) between two processes
 - * CPU to CPU transfer
 - * GPU to GPU transfer staged through CPU memory
 - * GPU to GPU tranfer through GPUDirect RDMA/NVLink
 - CUDA-aware MPI run config parameters: https://www.open-mpi.org/faq/?category=runcuda



Ping-pong Bandwidth





Summary

- Optimizations were done in Cray MPI to improve pt2pt and collective latency on KNL
- Further tuning is possible through the environment variables
- MPI on ThetaGPU
 - Recommend using the on-node NVLink communication for efficient data transfers

References:

- Cray XC series Network: https://www.cray.com/sites/default/files/resources/CrayXCNetwork.pdf
- MPI 3.1 Standard: https://www.mpi-forum.org/docs/mpi-3.1/mpi31-report.pdf
- Cray MPI for KNL: https://www.alcf.anl.gov/files/Chunduri_MPI_Theta.pdf (May 18 workshop slightly basic version than this talk)
- MPI benchmarking on Theta: https://cug.org/proceedings/cug2018_proceedings/includes/files/pap131s2-file1.pdf
- Advanced MPI Programming Tutorial at SC17, November 2017 (https://www.mcs.anl.gov/https://www.mcs.anl.gov/https://www.mcs.anl.gov/
- Low-overhead MPI profiling tool (Autoperf): https://www.alcf.anl.gov/user-guides/automatic-performance-collection-autoperf
- Run-to-run Variability: https://dl.acm.org/citation.cfm?id=3126908.3126926
- LDMS: https://github.com/ovis-hpc/ovis/tree/master/ldms



Hands-on Session

- Few sample codes are provided here at https://xgitlab.cels.anl.gov/alcf/training/tree/mpi/ProgrammingModels/MPI/Theta
- These are also accessible at /projects/Comp_Perf_Workshop/examples/training/ProgrammingModels/MPI/Theta
- The nonblocking_coll and nonblocking_p2p have sample run scripts to submit jobs
- Feel free to experiment with using the environment variables with your own application codes
- Some examples to try out
 - Potential performance optimization (Huge pages, Routing mode change, Hardware collective offload, Core specialization etc.)
 - Functionality specific (nonblocking collectives, MPI+X etc.)